

History

- Began in 1990 with a project for consumer electronic device controllers.
- 1991: Programming language "Oak" for this project.
- 1992: Separate company (FirstPerson, Inc.) started for this project, but didn't thrive.
- 1994: Sun began adapting Oak for the Internet.
- April 1995: Initial release (including HotJava), and rename to Java.
- May 1995: Netscape agrees to incorporate Java with Navigator.
- April 1996: Microsoft announces incorporation of Java into Windows 95.
- Following months: Lots and lots of people announce Java projects and support.

3

4

What makes Java so popular?

Java Overview

Java Overview

- Platform independence
- Pure (although slightly incomplete) object oriented design
- Secure for running untrusted code





Problem 1: Unchecked pointer arithmetic treats memory as one big block.

Example:

```
int var1;
int *ptr = &var1;
int var2 = *(ptr + 100); // This is legal but nonsensical!
```

Solution 1: No pointers in Java! *Reference* types in the Java VM, but cannot be mixed with integers (*no* pointer arithmetic), so at compile time can verify that they will always point to something reasonable.

Related problem/solution: All array bounds are checked before indexing is performed.

Problem 2: Stack problems/overwriting critical info.

9

10

Example: (Intel x86 assembly language)

mysubr proc

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pop ax mov ax, 0 push ax ret

mysubr endp

Solution 2: Stack operations are restricted to be predictable and bounded to remain in the appropriate region. Somewhat restrictive, but loss is efficiency not capability.

Not allowed: (Intel x86 assembly language)

mov bx, 10 mov cx, 0 prloop: xor dx, dx div bx push dx inc cx cmp ax, 0 jnz prloop

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Example:		
Object obj; // / int var; cout << obj.data; cout << var;	Assume no construc // Unpredictabl // Unpredictabl	tor that initializes val e value! e value!
Solution 3: All c default value. Loca run-time.	ass fields get an al variable initiali	initial, predictable izations verified at
Java Overview		1
Problem 4: Men (pointer/reference i anyway). Example:	ory allocation p not pointing to a	roblems valid object but used
char *str; strcpy(str, "Hell	o"); // What do	es str point to?!?
Solution 4: All d time.	ereference opera	tors checked at run

Problem 5: Memory deallocation problems (memory not freed properly or used after freeing).

Example 1 (basic problem):

```
Object *obj = new Object;
// ... use *obj
delete obj;
// ... more stuff
obj->data = 1;
```

Example 2 (more subtle):

```
Node *curr = first;
while (curr != 0) {
    delete curr;
    curr = curr->next; // *curr no longer there!
}
```

Example 3 (losing allocated memory):

Node *curr = new Node; curr = first; // What happened to new node?

Solution 5: No explicit deallocation available. All memory management is through automatic garbage collection, insuring consistency.

13

14

Java Overview

Problem 6: Types used inconsistently (at machine level or through unions in C/C++).

Example:

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union {
 int ival;
 float fval;
} u;
u.ival = 10;
cout << u.fval; // Not really a float value stored there!</pre>

Solution 6: Memory locations in Java VM have values with a *type*, so not just treated as strings of bits. Type consistency enforced at run-time (in Sun JDK this is done through load-time code validation).

Problem 7: Changing types/class signatures without proper recompilation of relevant modules causes "offset shift".

Example: Code compiled with this definition

```
class myclass {
public:
    int field1;
    float field2;
}:
```

may use memory offset 0 within the object to refer to field 1. If later changed to $\label{eq:constraint}$

```
class myclass {
public:
    float field2;
    int field1;
};
```

now previously compiled code does not work correctly because fields have moved!

Solution 7: Java bytecode (i.e., object code) contains no offsets (which would require pointer arithmetic!), only *name based* references resolved at run-time.

15

16

Java Overview

Problem 8: Overloading integers and booleans allows incorrect test conditions to slip by.

Example:

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if (x = 1)
 cout << "x is one!";</pre>

Solution 8: Addition of a separate boolean type, which is required for all test conditions. No casting between boolean and integer types.

Example: (Intel x86 machine code)			
0FFF:0100 B80001 0FFF:0103 E83D01 0FFF:0106 48 0FFF:0107 75FB 0FFF:0109 C3	MOV CALL DEC JNZ RET	AX,0100 0243 AX 0104	
Solution 9: All before the code is	jump targ executed	gets verified a	t loading time,
Java Overview			
Problem 10: For machine code allo deletions, etc.).	or untrust ws arbitra	ed (download ry OS calls (led) programs, file modificatior
Solution 10: <i>O</i> trusted libraries w manager class.	<i>nly</i> mach ith access	ine interface screened by	is through local a <i>security</i>

Final Note

These Java/JVM design decisions were made mostly to protect machine against *rogue code* or unpredictable behavior, but also excellent protection against *programmer error*!

- It's difficult to catch all the subtle machine-failure-type bugs in a C++ program.
- It's difficult for a subtle machine-failure-type bugs to get by testing in a Java program.

Of course, logic bugs are still the responsibility of the programmer! $% \left[\left({{{\mathbf{r}}_{i}}} \right) \right]$

19

Java Overview